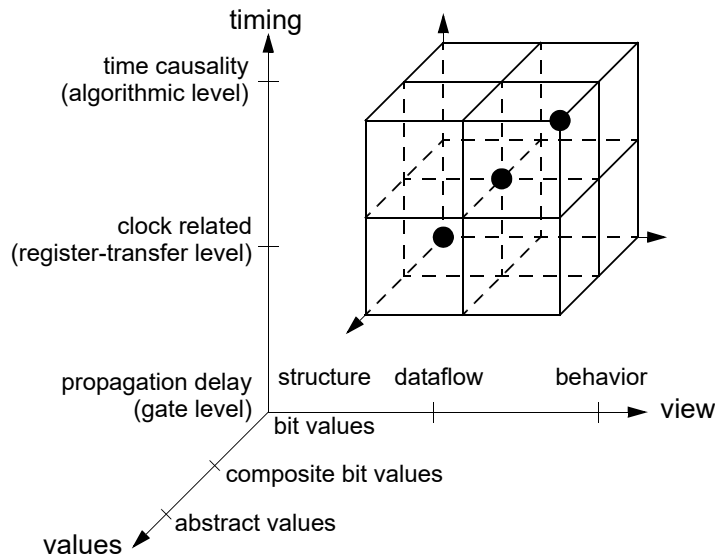


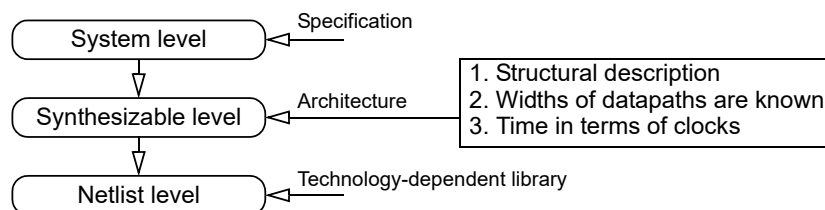


## HDL Design Cube



## Synthesis and VHDL

- VHDL LRM defines only simulation semantics of the language.
- LRM – Language Reference Manual



- **Synthesis restrictions:**
  - the lack of maturity of synthesis tools
  - the state-of-art in synthesis targets RTL synthesis only
  - certain VHDL features are simply not synthesizable
- The same applies to Verilog/SystemVerilog...



## Synthesis style

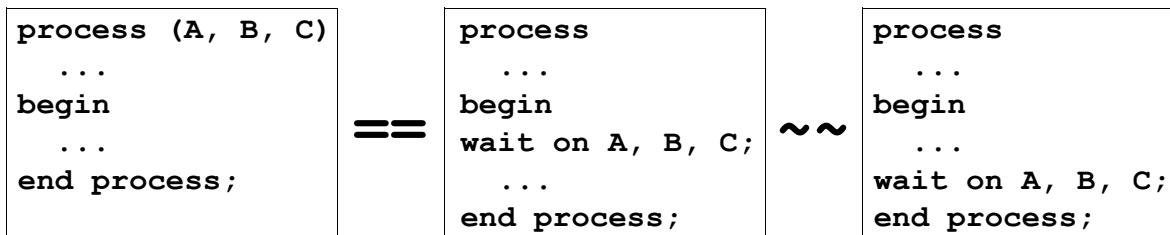
- Delay expressions (*after* clauses, *wait for* statements are ignored)
- Certain restrictions on the writing of process statement occur
- Only a few types are allowed
  - integer, enumerated, e.g., bit, bit\_vector, signed
- Type conversion and resolution functions are not interpreted
- Description is oriented towards synchronous styles with explicit clocks
- Types: enumeration, integer, one-dimensional array, record
 

```
type WORD is array (31 downto 0) of BIT;
type RAM is array (1023 downto 0) of WORD;
```
- In record, an item address becomes hardware coded
- !!! Time type is not supported !!!
- No explicit or default initialization
- Parenthesis in expressions have effect on HW generation
- Some arithmetic operations are supported partially only



## Sensitivity list

- Equivalent processes:



- Some synthesizers support only sensitivity list for combinational logic
- In case of single synchronization process there is no need to “remember” at which synchronization point it was stopped → such behavior does not imply memorization
- Process with multiple synchronization points, i.e. several wait states, infer memorization – FSM

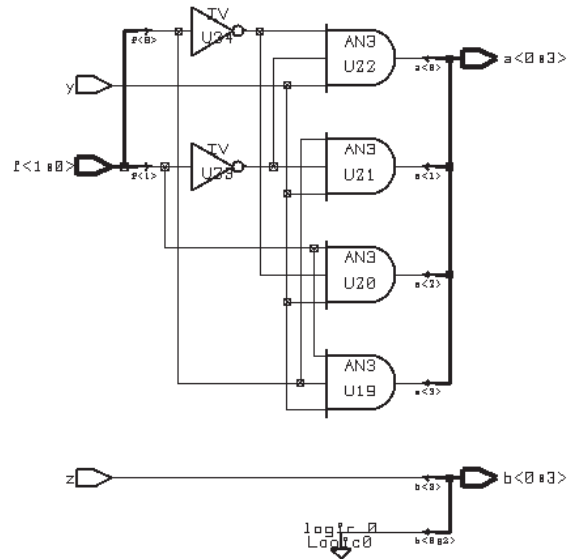


## Assignment statement synthesis

```

signal A,B: BIT_VECTOR(0 to 3);
signal I: INTEGER range 0 to 3;
signal Y,Z: BIT;
-- . . .
process ( I, Y, Z ) begin
  A<="0000";
  B<="0000";
  A(I)<=Y; -- Computable index
  B(3)<=Z; -- Constant index
end process;

```

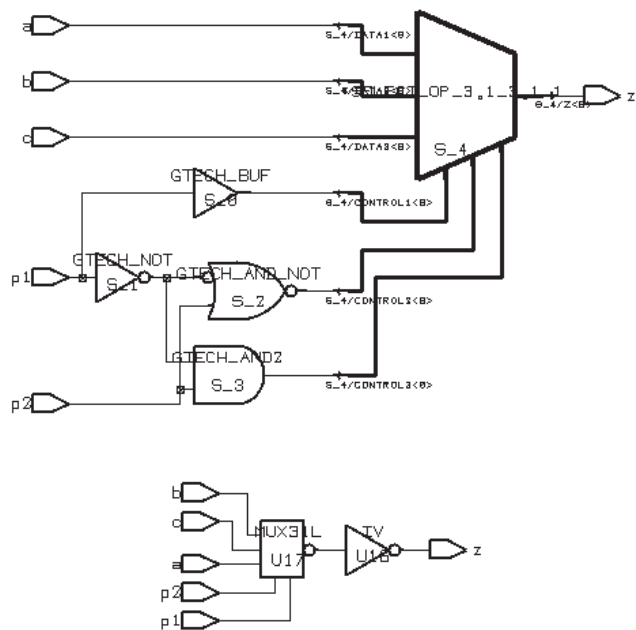


## if statement synthesis

```

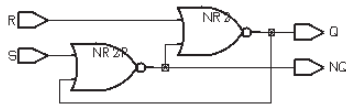
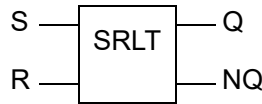
signal A,B,C,P1,P2,Z: BIT;
-- . . .
process (P1,P2,A,B,C) begin
  if (P1 = '1') then
    Z <= A;
  elsif (P2 = '0') then
    Z <= B;
  else
    Z <= C;
  end if;
end process;

```





## SR latch



```
use WORK.CHECK_PKG.all;
entity SRLT is
    port ( S, R:  in bit;
          Q, NQ: out bit );
begin
    NOT_AT_THE_SAME_TIME(S,R);
end SRLT;
```

```
architecture A1 of SRLT is
    signal LQ:  bit := '1';
    signal LNQ: bit := '0';
begin
    LNQ <= S nor LQ;
    LQ  <= R nor LNQ;
    Q   <= LQ;
    NQ  <= LNQ;
end A1;
```

- **NB! Asynchronous feed-back is temporarily cut by synthesizers...**



## Combinational circuit

- **A process is combinational, i.e. does not infer memorization, if:**
  - **the process has an explicit sensitivity list or contains a single synchronization point (waiting for changes on all input values); <sup>1)</sup>**
  - **no local variable declarations, or variables are assigned before being read;**
  - **all signals, which values are read, are part of the sensitivity list; <sup>2)</sup> and**
  - **all output signals are targets of signal assignments independent on the branch of the process, i.e. all signal assignments are covered by all conditional combinations.**

<sup>1)</sup> waiting on a clock signal, e.g., “ wait on clk until clk='1'; ”, implies buffered outputs (FF-s)

<sup>2)</sup> interpretation may differ from tool to tool

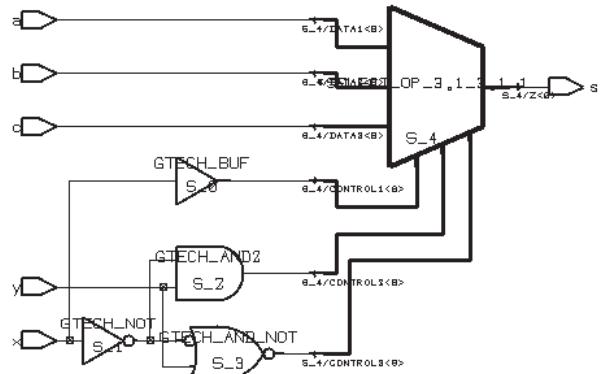
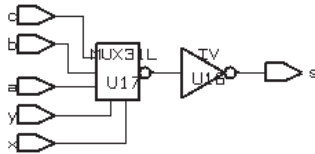


## Complex assignments

- No memory

`S <= A when X='1' else B when Y='1' else C;`

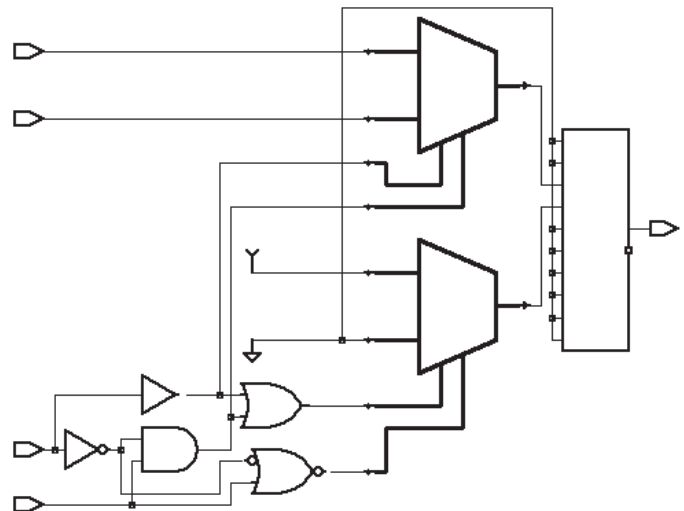
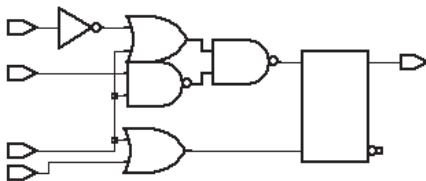
```
process (A, B, C, X, Y) begin
  if X='1' then S <= A;
  elsif Y='1' then S <= B;
  else S <= C;
  end if;
end process;
```



## Complex assignments (#2)

- Memory element generated!

```
process (A, B, X, Y) begin
  if X='1' then S <= A;
  elsif Y='1' then S <= B;
  end if;
end process;
```



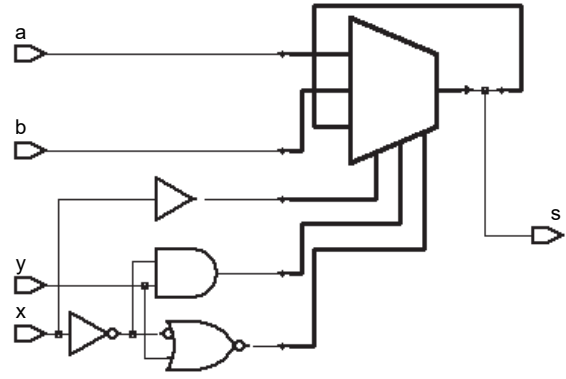
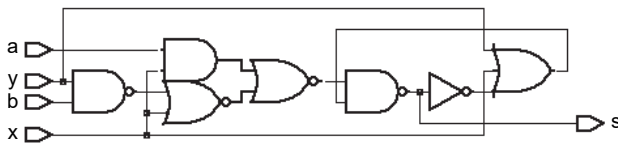


## Complex assignments (#3)

- Memory element generated!

```
S <= A when X='1' else B when Y='1' else S;
```

```
process (A, B, X, Y) begin
  if X='1' then S <= A;
  elsif Y='1' then S <= B;
  else S <= S;
  end if;
end process;
```

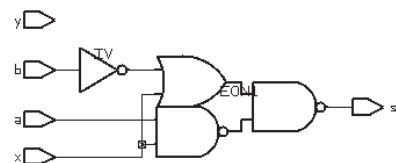
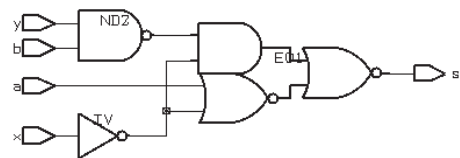


## Default values

- The default values inherited from type or subtype definitions
- The explicit initialization that is given when the object is declared
- A value assigned using a statement at the beginning of a process
- Only the last case is supported by synthesis tools!
- Usually, a part of the synthesizable code is devoted to *set/reset* constructions
- Default values can be used to guarantee that the signal always gets a new value

```
process (A, B, X, Y) begin
  S <= '0';
  if X='1' then S <= A;
  elsif Y='1' then S <= B;
  end if;
end process;
```

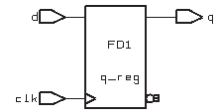
```
process (A, B, X, Y) begin
  S <= '-';
  if X='1' then S <= A;
  elsif Y='1' then S <= B;
  end if;
end process;
```



## Flip-flops

- **Process with the clock signal in the sensitivity list and explicit clock flank definition**  

```
process (CLK) begin
  if CLK='1' and CLK'event then  Q <= D;  end if;
end process;
```



- **Process with the clock signal and clock flank definition in the wait statement**  

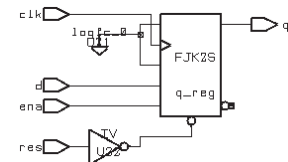
```
process begin
  wait on CLK until CLK='1';  Q <= D;
end process;
```

- **Concurrent assignment with the clock signal and clock flank definition**  

```
Q <= D when CLK='1' and CLK'event;
```

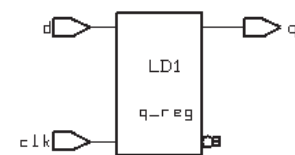
- **Asynchronous reset & synchronous enable**  

```
process (RES,CLK) begin
  if RES='1' then  Q <= '0';  -- asynchronous reset
  elsif CLK='1' and CLK'EVENT then
    if ENA='1' then  Q <= D;  end if;
  end if;
end process P4_FF;
```



## Latch vs. Flip-flop?

```
P1_L: process (CLK, D) begin
  if CLK='1' then  Q <= D;
  end if;
end process P1_L;
```



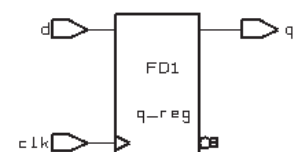
```
P2_FL: process (CLK) begin
  if CLK='1' then  Q<=D;
  end if;
end process P2_FL;
```

- **Simulation OK but not synthesis!**  

```
-- Warning: Variable 'd' is being read
-- in routine .. line .. in file '..',
-- but is not in the process sensitivity
-- list of the block which begins
-- there. (HDL-179)
```

- **Result - latch**

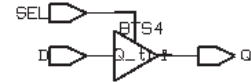
```
P1_FF: process (CLK) begin
  if CLK='1' and
    CLK'event then  Q<=D;
  end if;
end process P1_FF;
```





## Synthesis rules

- Guidelines in priority order:
  - the target signal(s) will be synthesized as flip-flops when there is a signal edge expression, e.g. CLK'event and CLK='1', in the process
  - usually, only one edge expression is allowed per process
    - different processes can have different clocks (tool depending)
  - the target signal will infer three-state buffer(s) when it can be assigned a value 'Z'
    - example: Q <= D when SEL='1' else 'Z';
  - the target signal will infer a latch (latches) when the target signal is not assigned with a value in every conditional branch, and the edge expression is missing
  - a combinational circuit will be synthesized otherwise
- It is a good practice to isolate flip-flops, latches and three-state buffers inferences to ensure design correctness



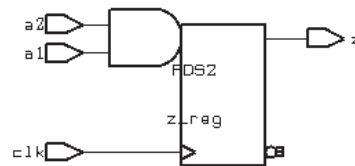
## Signal versus variable

- The hardware resulting from synthesis of variables or signals differs: either nothing, wires, or memory elements

```

signal A1, A2: BIT;
-- . . .
process (CLOCK)
  variable A3: BIT;
begin
  if CLOCK='1' and CLOCK'event then
    A3 := A1 and A2;
    Z <= A3;
  end if;
end process;

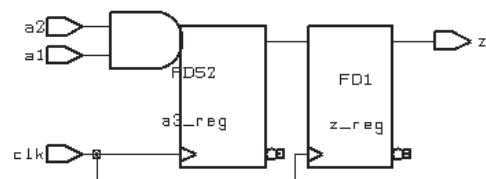
```



```

signal A1, A2, A3: BIT;
-- . . .
process (CLOCK)
begin
  if CLOCK='1' and CLOCK'event then
    A3 <= A1 and A2;    -- Next delta-cycle!!
    Z <= A3;
  end if;
end process;

```



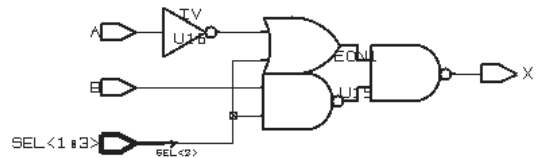


# Arithmetics

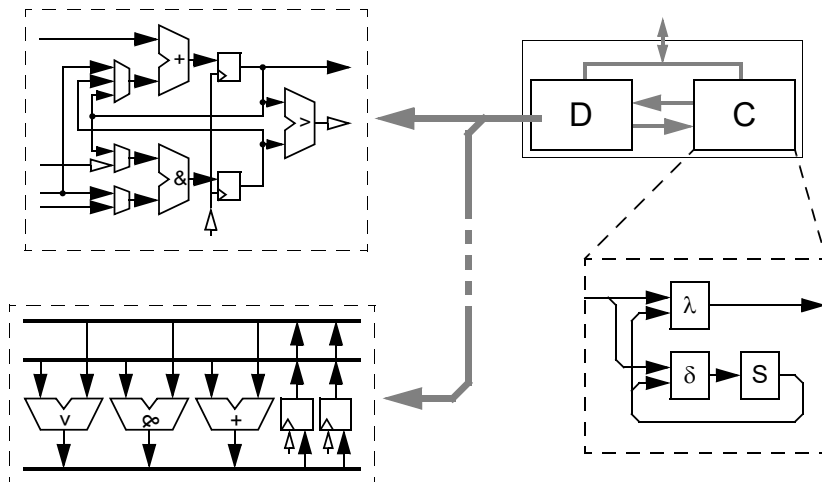
- **Overloaded arithmetic operations:**  

```
library IEEE;
use IEEE.std_logic_1164.all;
use IEEE.std_logic_arith.all;
```
- **Sources of the named packages are in the directory: \$SYNOPSIS/packages/IEEE/src**
  - \$SYNOPSIS/packages is the root directory for all Synopsys packages
- **Be careful with ‘\*’, ‘/’, ‘\*\*’ - extremely chip area consuming**
  - Safe in some special cases - multiplication by power of two
- **Use parenthesis to group a set of gates**
- **Don't care values and synthesis**

```
process (A, B, SEL) begin
  case SEL is
    when "001" => X <= A;
    when "010" => X <= B;
    when others => X <= '-';
  end case;
end process;
```



# Data-part & control-part

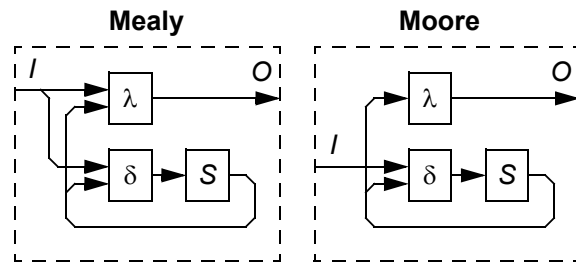


- **one unit – one process**
- **functional units – combinational processes** [all inputs in the sensitivity list]
- **storage units – clocked processes** [activation at clock edge]



## FSM in VHDL

- **FSM:**  $M = (S, I, O, \delta, \lambda)$ 
  - **S:** set of states
  - **I:** set on inputs
  - **O:** set of outputs
  - $\delta$ : transition function -  $\delta: S \times I \rightarrow S$
  - $\lambda$ : output function -  $\lambda: S \times I \rightarrow O$
- **How many processes?**
  - Process per block
- **Three processes**
  - (1) transition function, (2) output function, (3) state register
- **Two processes**
  - (1) merged transition and output functions, (2) state register [Mealy]
- **One process**
  - buffered outputs! [Moore]



## FSM as a single process

- Note that all signals assigned in the process will have flip-flops!
 

```

-- RESET is the asynchronous reset, CLK is the clock
-- STATE is a variable (or signal) memorizing the current state
process (RESET,CLK)
begin
    if RESET='1' then          -- asynchronous reset
        STATE <= S_INIT;
    elsif CLK='1' and CLK'EVENT then
        case STATE is
            when S_INIT => if I0='1' then STATE <= S5; end if;
            when ... => ...
        end case;
    end if;
end process;

```



## Three process FSM

- storage elements, transition function & output function

```

architecture B of FSM is
    type TYPE_STATE is (S_INIT,S1,...Sn);
    signal CURRENT_STATE, NEXT_STATE : TYPE_STATE;
begin
    P_STATE: process begin -- sequential process / storage elements
        wait until CLK'EVENT and CLK='1';
        if RESET='1' then CURRENT_STATE <= S_INIT;
        else
            CURRENT_STATE <= NEXT_STATE; end if;
    end process P_STATE;
    P_NEXT_STATE: process (I0, ..., CURRENT_STATE) begin -- next state function
        NEXT_STATE <= CURRENT_STATE;
        case CURRENT_STATE is
            when S_INIT => if I0='1' then NEXT_STATE <= S5; end if;
            when ... => ...
        end case;
    end process P_NEXT_STATE;
    P_OUTPUTS: process (CURRENT_STATE) begin -- output function
        case is
            when S_INIT => O <= '0';
            when ... => ...
        end case;
    end process P_OUTPUTS;
end B;

```



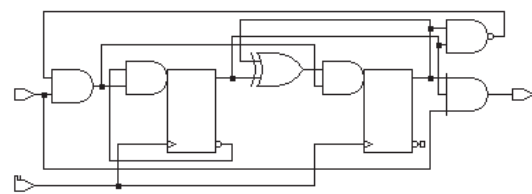
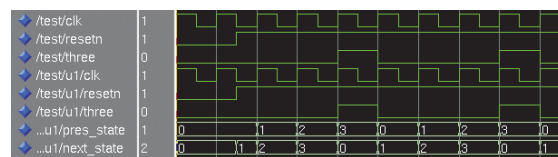
## FSM #2 – description styles & synthesis

### Two processes (modulo-4 counter)

```

library IEEE; use IEEE.std_logic_1164.all;
entity counter03 is
    port ( clk: in bit;
          resetn: in std_logic;
          three: out std_logic );
end entity counter03;
architecture fsm2 of counter03 is
    subtype state_type is integer range 0 to 3;
    signal pres_state, next_state: state_type := 0;
begin
    process (clk) begin -- State memory
        if clk'event and clk = '1' then
            pres_state <= next_state;
        end if;
    end process;
    -- Next state & output functions
    process (resetn, pres_state) begin
        three <= '0';
        if resetn='0' then next_state <= 0;
        else
            case pres_state is
                when 0 to 2 => next_state <= pres_state + 1;
                when 3 => next_state <= 0; three <= '1';
            end case;
        end if;
    end process;
end architecture fsm2;

```



22 gates / 3.70 ns

## FSM #2 – description styles & synthesis

### Three processes (modulo-4 counter)

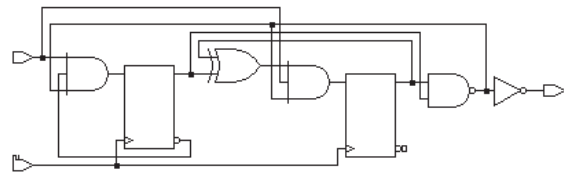
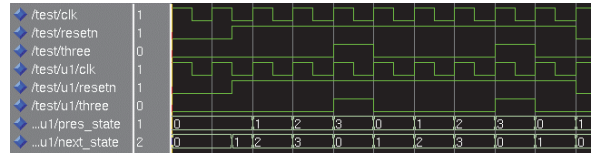
```

library IEEE; use IEEE.std_logic_1164.all;
architecture fsm3 of counter03 is
  subtype state_type is integer range 0 to 3;
  signal pres_state, next_state: state_type := 0;
begin
  process (clk) begin -- State memory
    if clk'event and clk = '1' then
      pres_state <= next_state;
    end if;
  end process;

  -- Next state function
  process (resetn, pres_state) begin
    if resetn='0' then next_state <= 0;
    else
      if pres_state=3 then next_state <= 0;
      else next_state <= pres_state + 1;
      end if;
    end if;
  end process;

  -- Output function
  process (resetn, pres_state) begin
    if pres_state=3 then three <= '1';
    else three <= '0';
    end if;
  end process;
end architecture fsm3;

```



23 gates / 4.36 ns

## FSM #2 – description styles & synthesis

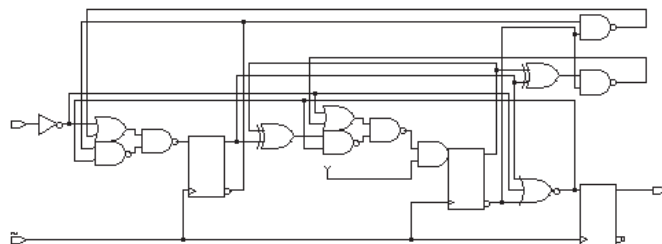
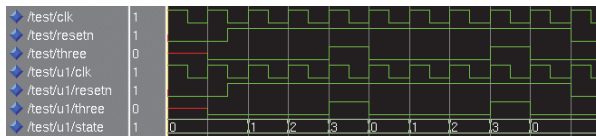
### One process (modulo-4 counter)

```

library IEEE; use IEEE.std_logic_1164.all;
architecture fsm1 of counter03 is
  subtype state_type is integer range 0 to 3;
  signal state: state_type := 0;
begin
  process (clk) begin
    if clk'event and clk = '1' then
      three <= '0';
      if resetn='0' then state <= 0;
      else
        case state is
          when 0 | 1 => state <= state + 1;
          when 2 => state <= state + 1; three <= '1';
          when 3 => state <= 0;
        end case;
      end if;
    end if;
  end process;
end architecture fsm1;

// Another version to build the process
process begin
  wait on clk until clk='1';
  three <= '0';
  if resetn='0' then state <= 0;
  else
    -- and so on...

```



38 gates / 5.68 ns



## Using generics

```
entity AND_N is
    generic (N: POSITIVE);
    port    (DIN: in BIT_VECTOR(1 to N); R: out BIT);
end AND_N;

architecture A1 of AND_N is
    signal INTER: BIT_VECTOR(1 to N);
begin
    INTER(1) <= DIN(1);
    L: for I in 1 to N-1 generate
        INTER(I+1) <= DIN(I+1) and INTER(I);
    end generate;
    R <= INTER(N);
end A1;

C1: AND_N generic map (N=>12) port map(IN_DATA, OUT_DATA);
```



## for-loop versus while-loop?

- **May be tool dependent!**
  - Design Compiler (Synopsys): *for* - parallel, *while* - sequential
  - ISE (Xilinx): *for* / *while* - both parallel
  - Leonardo (Mentor Graphics): depending on the timing constructs
  
- **for-loop**
  - parallel implementation
  - no timing control (wait) in the loop body

```
for i in 0 to 7 loop
    x(i) := a(i) and b(i);
end loop;
```
  
- **while-loop**
  - sequential implementation
  - timing control (wait) required in the loop body

```
i := 0;
while i<7 loop
    data(i) := in_port;
    wait on clk until clk='1';
    i := i + 1;
end loop;
```



## Multiple wait statements

- VHDL semantics must be preserved
  - different interpretations possible
- Distributing operations over multiple clock steps

- **Algorithm**

- Inputs: a, b, c, d
- Output: x
- Coefficients: c1, c2
- $x = a + b \cdot c1 + c \cdot c2 + d$
- Timing constraint - 3 clock periods

```

process
  variable av, bv, cv, dv: ...;
begin
  av:=a; bv:=b; cv:=c; dv:=d;
  wait on clk until clk='1';
  wait on clk until clk='1';
  x <= av + bv * c1 + cv * c2 + dv;
  wait on clk until clk='1';
end process;
    
```



## Multiple wait statements

- Behavioral interpretation may lead to an unoptimal solution

```

process
  variable av, bv, cv, dv: ...;
begin
  av:=a; bv:=b; cv:=c; dv:=d;
  wait on clk until clk='1';
  wait on clk until clk='1';
  x <= av + bv * c1 + cv * c2 + dv;
  wait on clk until clk='1';
end process;
    
```

2 multipliers & 3 adders

```

process
  variable av, bv, cv, dv: ...;
  variable r1, r2: ...;
begin
  av:=a; bv:=b; cv:=c; dv:=d;
  r1 := av + dv;    r2 := bv * c1;
  wait on clk until clk='1';
  r1 := r1 + r2;    r2 := cv * c2;
  wait on clk until clk='1';
  x <= r1 + r2;
  wait on clk until clk='1';
end process;
    
```

1 multiplier & 1 adder

*Behavioral Synthesis  
(High-Level Synthesis)*



## Inserting wait statements

- VHDL semantics preserved for inputs/outputs
- targeting as-fast-as-possible (AFAP) schedules

- **16-tap FIR filter**

- new input and output data at every rising flank of sys\_clk (sampling clock)
- internal clock can be added

- **How to implement loops?**

- 1st - in parallel (shift-register)
- 2nd - sequentially
  - multiply-and-accumulate (MAC)
  - ROM for coefficients

```

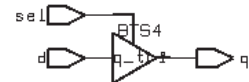
process
  variable sum: ...;
  variable buff: ...; -- array (0 to 15)
begin
  for i in 15 downto 1 loop
    buff(i) := buff(i-1);
  end loop;
  buff(0) := data_in;  sum:=0;
  for i in 0 to 15 loop
    sum := sum + buff(i) * coeff(i);
  end loop;
  x <= sum;
  wait on sys_clk until sys_clk='1';
end process;
    
```



## Verilog/SystemVerilog – synthesis rules

- **Guidelines in priority order:**

- the target signal(s) will be synthesized as flip-flops when there is a signal edge expression, e.g. “ @(posedge CLK) ”, in the behavioral statement
- only one edge expression is allowed per behavioral statement
  - different statements can have different clocks (tool depending)
- the target signal will infer three-state buffer(s) when it can be assigned a value 'Z'
  - example: q = sel == 1 ? d : 'bz';
- the target signal will infer a latch (latches) when the target signal is not assigned with a value in every conditional branch, and the edge expression is missing
- a combinational circuit will be synthesized otherwise



- It is a good practice to isolate flip-flops, latches and three-state buffers inferences to ensure design correctness



## Combinational circuit

- A process is combinational, i.e. does not infer memorization, if:
  - the behavioral statement has a sensitivity list in the beginning (waiting for changes on all input values); <sup>1)</sup>
  - signals are assigned before being read;
  - all signals, which values are read, are part of the sensitivity list; <sup>2)</sup> and
  - all output signals are targets of signal assignments independent on the branch of the process, i.e. all signal assignments are covered by all conditional combinations.

1) waiting on a clock signal, e.g., “ @(posedge clk) ”, implies buffered outputs (FF-s)

2) interpretation may differ from tool to tool

- SystemVerilog has three new *always* constructs
  - always\_comb – explicit combinational circuit
  - always\_latch – explicit latch
  - always\_ff – explicit flip-flop



## Sensitivity list

- Equivalent statements:

```
always
  @(a or b or c or x or y)
begin
  if (x==1)      s=a;
  else if (y==1) s=b;
  else           s=c;
end
```

==

```
always begin
  @(a or b or c or x or y);
  if (x==1)      s=a;
  else if (y==1) s=b;
  else           s=c;
end
```

- In case of single synchronization process there is no need to “remember” at which synchronization point it was stopped → such behavior does not imply memorization





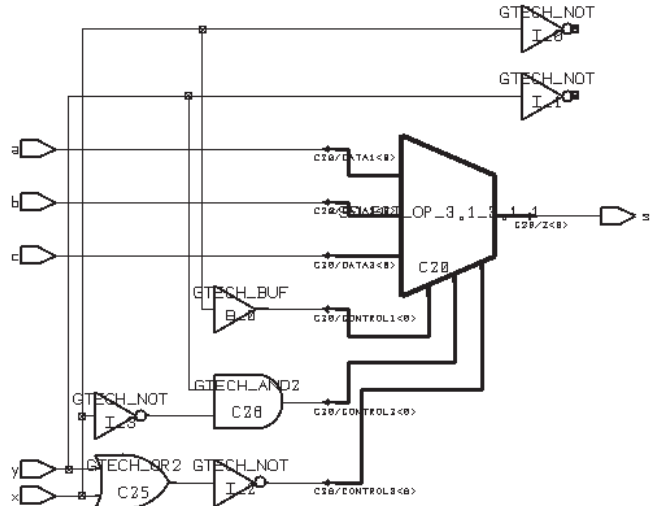
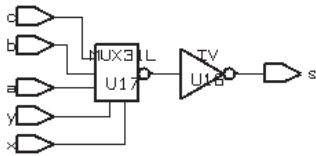
## Complex assignments

- No memory

```
assign s = x==1 ? a : y==1 ? b : c;
```

```
always
  @(a or b or c or x or y)
  if (x==1)      s=a;
  else if (y==1) s=b;
  else           s=c;
```

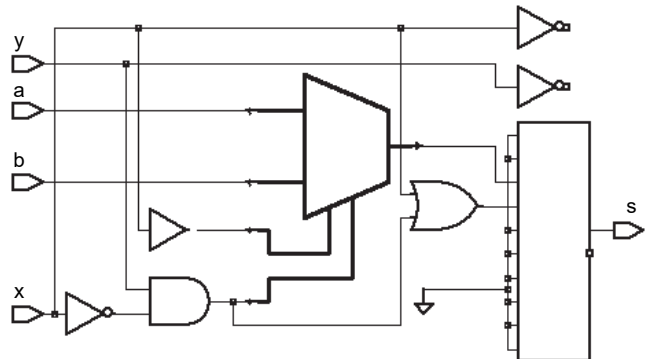
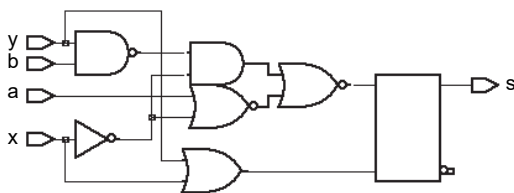
```
always_comb
  if (x==1)      s=a;
  else if (y==1) s=b;
  else           s=c;
```



## Complex assignments (#2)

- Memory element generated!

```
always begin
  @(a or b or x or y);
  if (x==1)      s=a;
  else if (y==1) s=b;
end
```

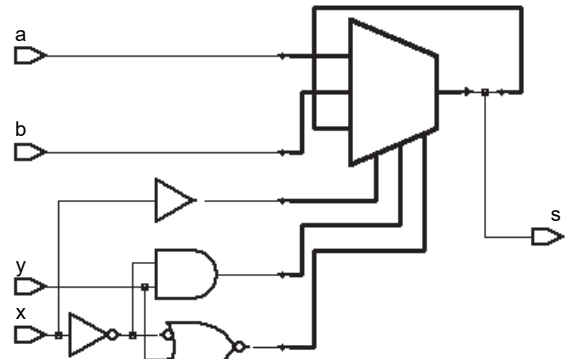
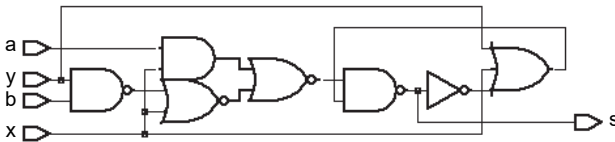




## Complex assignments (#3)

- Memory element generated!

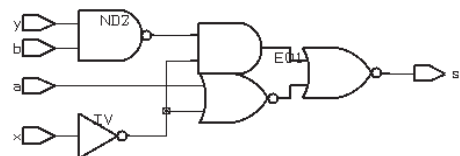
```
always begin
  @(a or b or x or y);
  if (x==1)      s=a;
  else if (y==1) s=b;
  else          s=s;
end
```



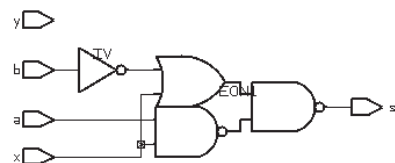
## Default values

- The default values inherited from type or subtype definitions
- The explicit initialization that is given when the object is declared
- A value assigned using a statement at the beginning of a process
- Only the last case is supported by synthesis tools!
- Usually, a part of the synthesizable code is devoted to *set/reset* constructions
- Default values can be used to guarantee that the signal always gets a new value

```
always begin
  @(a or b or x or y);
  s=0;
  if (x==1)      s=a;
  else if (y==1) s=b;
end
```



```
always begin
  @(a or b or x or y);
  s='bx';
  if (x==1)      s=a;
  else if (y==1) s=b;
end
```





## Latches & flip-flops

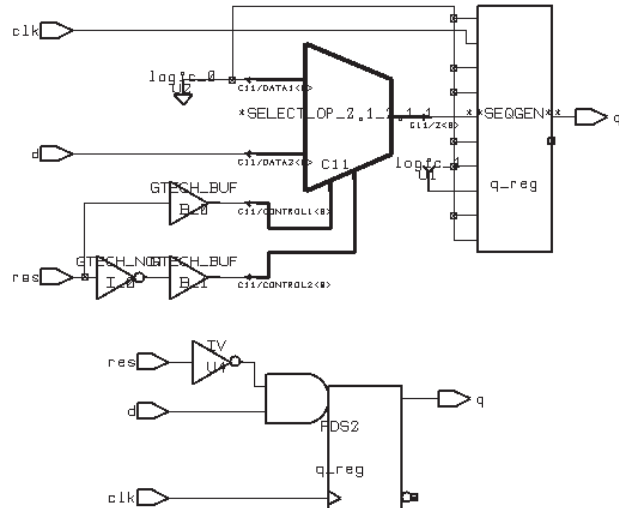
```
always @(clk or d) if (clk) q = d;    /* Latch */
```

```
always @(posedge clk) q = d;        /* Flip-flop */
always @(posedge clk) q <= d;
```

- synchronous reset

```
always @(posedge clk)
  if (res==1) q = 0;
  else      q = d;
```

```
always begin
  @(posedge clk);
  if (res==1) q = 0;
  else      q = d;
end
```

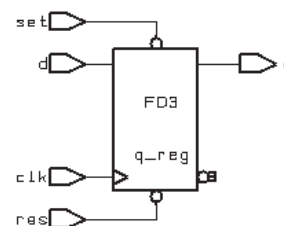


## Flip-flops

- asynchronous reset

```
always
  @(posedge res or
   posedge clk)
  if (res==1) q = 0;
  else      q = d;
```

```
always
  @(negedge res or
   negedge set or
   posedge clk)
  if (res==0)      q = 0;
  else if (set==0) q = 1;
  else            q = d;
```

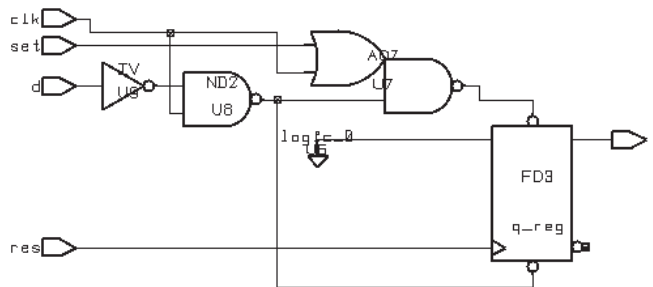




## Flip-flops

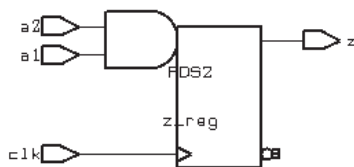
- asynchronous reset - the order of signals!

```
always @(posedge res or posedge set or posedge clk)
  if (clk==1)      q = d;
  else if (set==1) q = 1;
  else             q = 0;
```

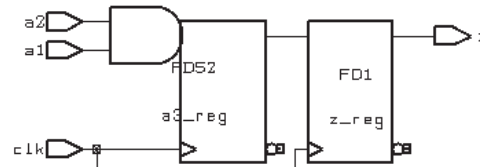


## Blocking versus non-blocking

```
module sig_var_b (clk, a1, a2, z);
  input clk, a1, a2;
  output z; reg a3;
  always @(posedge clk) begin
    a3 = a1 & a2;
    z <= a3;
  end
endmodule // sig_var_b
```



```
module sig_var_n (clk, a1, a2, z);
  input clk, a1, a2;
  output z; reg a3;
  always @(posedge clk) begin
    a3 <= a1 & a2;
    z <= a3;
  end
endmodule // sig_var_n
```



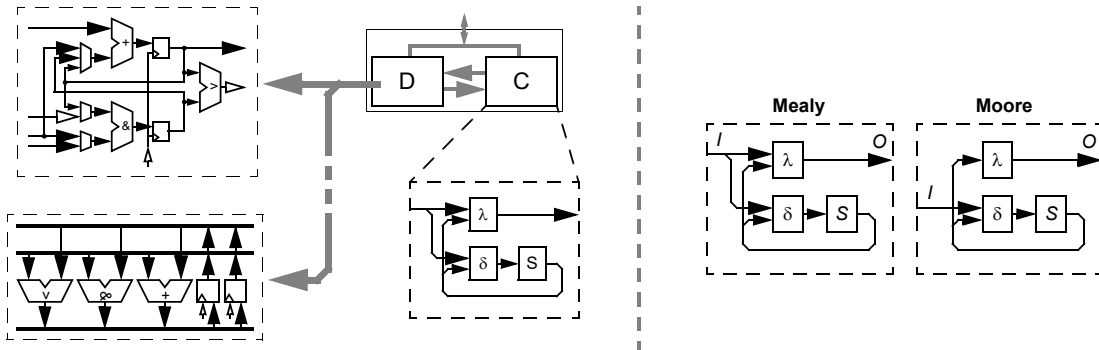
## Compare – signal versus variable in VHDL

```
process (CLK)
  variable A3 : BIT;
begin
  if CLK'event and CLK='1' then
    A3 := A1 and A2;
    Z <= A3;
  end if;
end process;
```

```
signal A3 : BIT;
-- ...
process (CLK) begin
  if CLK'event and CLK='1' then
    A3 <= A1 and A2;
    Z <= A3;
  end if;
end process;
```



## Data-part, control-part & FSM



- **one unit – one process**
  - functional units – combinational processes [all inputs in the sensitivity list]
  - storage units – clocked processes [activation at clock edge]
- **FSM:  $M = (S, I, O, \delta, \lambda)$  – process per block**
- **Three processes – (1) transition function, (2) output function, (3) state register**
- **Two processes – (1) merged transition and output functions, (2) state register [Mealy]**
- **One process – buffered outputs! [Moore]**

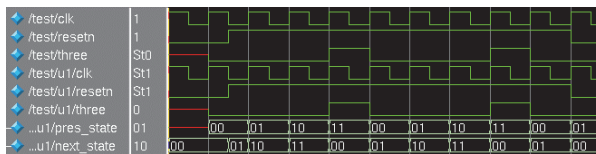


## FSM - description styles

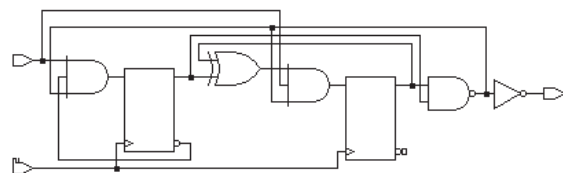
*Three processes* (modulo-4 counter)

```

module counter03 ( clk, resetn, three );
input clk, resetn;
output three;   reg three;
reg [1:0] pres_state, next_state;
always @(posedge clk) // State memory
    pres_state <= next_state;
// Next state function
always @(resetn or pres_state) begin
    if (resetn==0) next_state = 0;
    else case (pres_state)
        0, 1, 2: next_state = pres_state + 1;
        3:      next_state = 0;
    endcase
end
always @(pres_state) // Output function
    if (pres_state==3) three = 1;
    else                 three = 0;
endmodule
    
```



23 gates / 4.36 ns



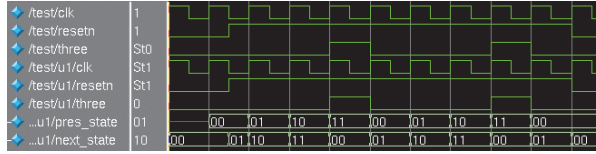
# FSM - description styles

## Two processes (modulo-4 counter)

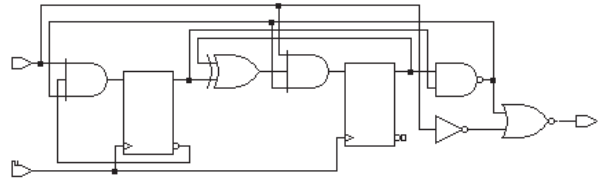
```

module counter03 ( clk, resetn, three );
input clk, resetn;
output three; reg three;
reg [1:0] pres_state, next_state;
always @(posedge clk) // State memory
pres_state = next_state;
// Next state & output functions
always @(resetn or pres_state) begin
three = 0;
if (resetn==0) next_state = 0;
else
case (pres_state)
0, 1, 2: next_state = pres_state + 1;
3: begin next_state = 0; three = 1; end
endcase
end
endmodule

```



24 gates / 4.36 ns



# FSM - description styles

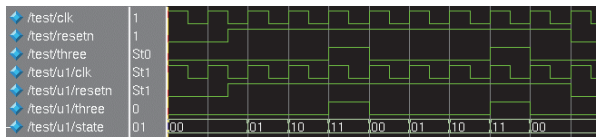
## One process (modulo-4 counter)

```

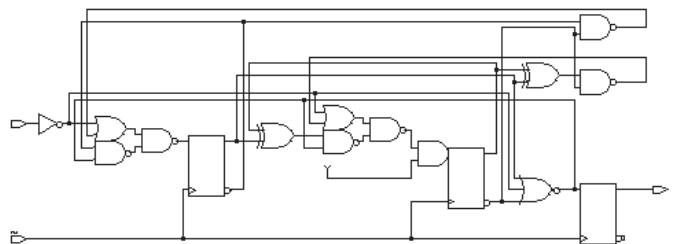
module counter03 ( clk, resetn, three );
input clk, resetn;
output three; reg three;
reg [1:0] state;
always @(posedge clk) begin
three = 0;
if (resetn==0) state = 0;
else case (state)
0, 1: state = state + 1;
2: begin state = state + 1; three = 1; end
3: state = 0;
endcase
end
endmodule

// Another version
// to begin the always block
always begin @(posedge clk);
three = 0; // and so on...

```



38 gates / 5.68 ns



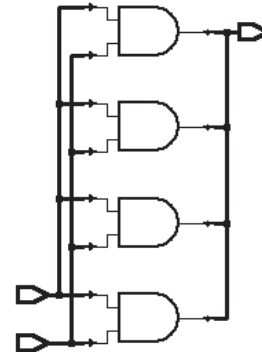


## for-loop versus while-loop?

- Is tool dependent!
  - Design Compiler (Synopsys) & ISE (Xilinx): *for* - parallel, *while* - parallel
  - No multiple waits!

```
always @(a or b) begin
  for (i=0;i<4;i=i+1)
    x[i] = a[i] & b[i];
end
```

```
always @(a or b) begin
  i = 0;
  while (i<4) begin
    x[i] = a[i] & b[i];
    i = i + 1;
  end
end
```

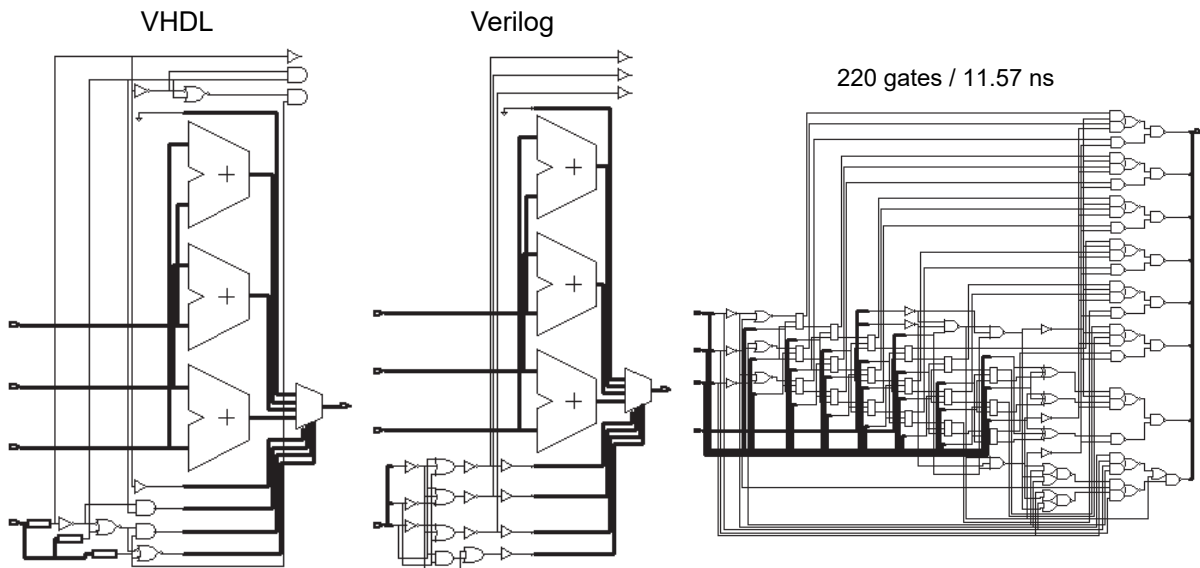


## Behavioral RTL vs. "pure" RTL

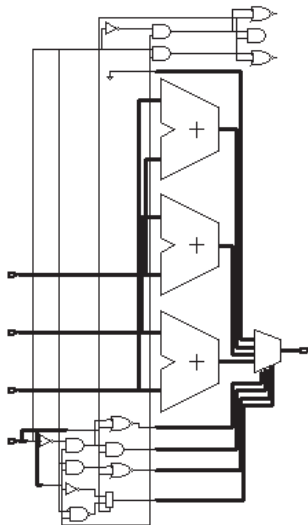
```
library IEEE;
use IEEE.std_logic_1164.all;
use IEEE.std_logic_arith.all;
entity test is
  port ( a, b, c: in unsigned(7 downto 0);
        x: in unsigned(2 downto 0);
        o: out unsigned(7 downto 0) );
end test;
architecture bhv of test is begin
process (a, b, c, x)
  constant x2: unsigned(2 downto 0):="010";
  constant x3: unsigned(2 downto 0):="011";
  constant x6: unsigned(2 downto 0):="110";
begin
  if x=x2 then o <= a+b;
  elsif x=x3 then o <= a+c;
  elsif x=x6 then o <= b+c;
  else o <= (others=>'0');
  end if;
end process;
end architecture bhv;
```

```
module test (a, b, c, x, o);
  input [7:0] a, b, c;
  input [2:0] x;
  output [7:0] o; reg [7:0] o;
  always @(a or b or c or x)
    if (x==2) o <= a+b;
    else if (x==3) o <= a+c;
    else if (x==6) o <= b+c;
    else o <= 0;
endmodule // test
```

## Behavioral RTL vs. "pure" RTL



## Behavioral RTL vs. "pure" RTL

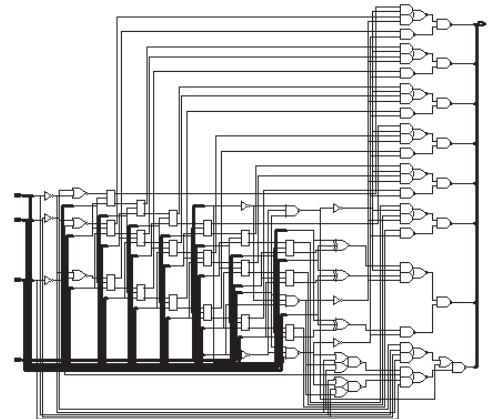


```

architecture bhv2 of test is begin
process (a, b, c, x) begin
  case x is
    when "010" => o <= a+b;
    when "011" => o <= a+c;
    when "110" => o <= b+c;
    when others => o <= (others=>'0');
  end case;
end process;
end architecture bhv2;

```

220 gates / 11.57 ns







## Behavioral RTL vs. "pure" RTL

```

architecture rtl of test is
    signal a1, a2: unsigned(7 downto 0);
    signal dc: unsigned(1 downto 0);
begin
    dec: process (x) begin
        case x is
            when "010" => dc <= "01";
            when "011" => dc <= "10";
            when "110" => dc <= "11";
            when others => dc <= "00";
        end case;
    end process dec;
    m1: process (a, b, dc) begin
        case dc is
            when "01" => a1 <= a;
            when "10" => a1 <= a;
            when "11" => a1 <= b;
            when others => a1 <= (others=>'0');
        end case;
    end process m1;
    m2: process (b, c, dc) begin
        -- ...
    end process m2;
    o <= a1 + a2;
end architecture rtl;

```

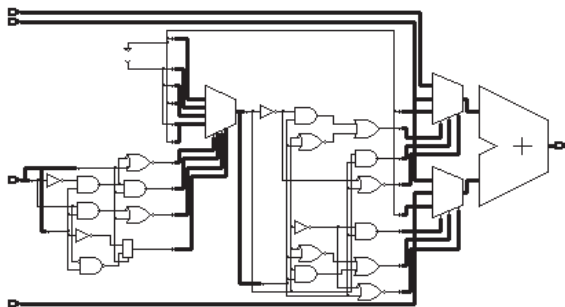
```

module test (a, b, c, x, o);
    input [7:0] a, b, c;
    input [2:0] x;
    output [7:0] o;
    reg [7:0] a1, a2;
    reg [2:0] dc;
    always @(x)
        if (x==2) dc = 1;
        else if (x==3) dc = 2;
        else if (x==6) dc = 3;
        else dc = 0;
    always @(a or b or dc)
        if (dc==1) a1 = a;
        else if (dc==2) a1 = a;
        else if (dc==3) a1 = b;
        else a1 = 0;
    always @(b or c or dc)
        if (dc==1) a2 = b;
        else if (dc==2) a2 = c;
        else if (dc==3) a2 = c;
        else a2 = 0;
    assign o = a1+a2;
endmodule // test

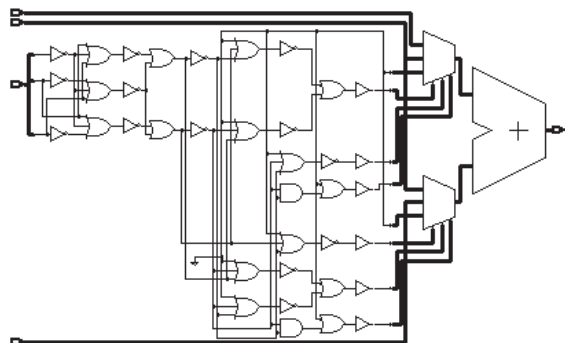
```



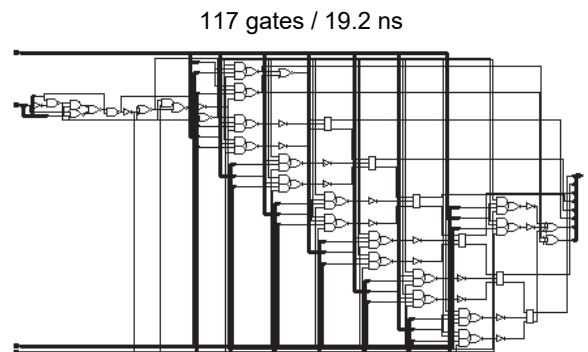
## Behavioral RTL vs. "pure" RTL



VHDL



Verilog





## Adder / Subtractor

```
library IEEE;
use IEEE.std_logic_1164.all;
use IEEE.std_logic_arith.all;
entity add_sub is
  port ( a, b: in unsigned(7 downto 0);
        x: in std_logic;
        o: out unsigned(7 downto 0) );
end add_sub;
architecture bhv of add_sub is begin
process ( a, b, x) begin
  if x='0' then o <= a+b;
  else o <= a-b; end if;
end process;
end architecture bhv;
```

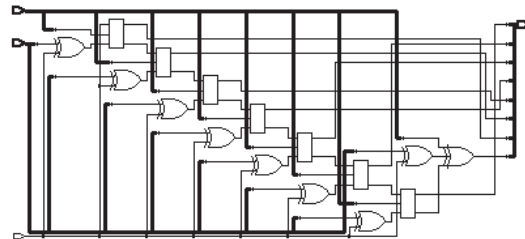
```
module add_sub (a, b, x, o);
  input [7:0] a, b;
  input x;
  output [7:0] o;
  assign o = x==0 ? a+b : a-b;
endmodule // add_sub
```

145 gates / 11.64 ns

```
architecture dfl of test5 is
  signal a1, b1, o1: unsigned(8 downto 0);
begin
  a1 <= a & '1';
  b1 <= b & '0' when x='0' else
    unsigned(not std_logic_vector(b)) &
    '1';
  o1 <= a1+b1;
  o <= o1(8 downto 1);
end architecture dfl;
```

```
/* ... */
assign {o,t} = {a,1'b1} +
  ( x==0 ? {b,1'b0} : {~b,1'b1} );
```

87 gates / 12.45 ns

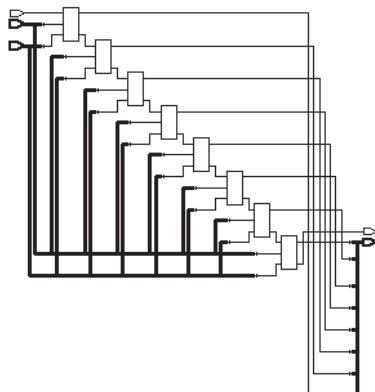


## Adders & Subtractors

```
a1 <= '0' & a & '1';
b1 <= '0' & b & ci;
o1 <= a1 + b1;
o <= o1(8 downto 1);
co <= o1(9);

assign {co,o,t} = {1'b0,a,1'b1} +
  {1'b0,b,ci};
```

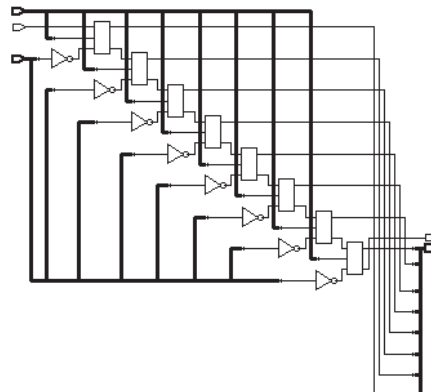
64 g. / 10.66 ns [ 60 g. / 10.08 ns w/o ci/co ]



```
a1 <= '0' & a & '1';
b1 <= '0' &
  unsigned(not std_logic_vector(b)) & ci;
o1 <= a1 + b1;
o <= o1(8 downto 1);
co <= o1(9);

assign {co,o,t} = {1'b0,a,1'b1} +
  {1'b0,~b,ci};
```

72 g. / 10.62 ns [ 66 g. / 10.35 ns w/o ci/co ]





## GCD (Greatest Common Divisor) example

- Specification ~ behavioral description
  - input/output timing fixed – handshaking signals & clock

```
process -- gcd-bhv.vhdl
    variable x, y: unsigned(15 downto 0);
begin
    -- Wait for the new input data
    wait on clk until clk='1' and rst='0';
    x := xi;    y := yi;    rdy <= '0';
    wait on clk until clk='1';
    -- Calculate
    while x /= y loop
        if x < y then y := y - x;
        else          x := x - y;    end if;
    end loop;
    -- Ready
    xo <= x;    rdy <= '1';
    wait on clk until clk='1';
end process;
```

### Problems

- inner loop not clocked
- complex wait statement
- ( - multiple wait statements )

### What to look for?

- different synthesis tools
- minimizing resources
- maximizing performance

### Target technologies – ASIC, FPGA

### VHDL code & testbenches

<http://mini.pld.ttu.ee/~lrv/gcd/gcd.html>



## GCD example – synthesizable code?

- Clocked behavioral style

```
process -- gcd-bhvc.vhdl
    variable x, y: unsigned(15 downto 0);
begin
    -- Wait for the new input data
    while rst = '1' loop
        wait on clk until clk='1';
    end loop;
    x := xi;    y := yi;    rdy <= '0';
    wait on clk until clk='1';
    -- Calculate
    while x /= y loop
        if x < y then y := y - x;
        else          x := x - y;    end if;
        wait on clk until clk='1';
    end loop;
    -- Ready
    xo <= x;    rdy <= '1';
    wait on clk until clk='1';
end process;
```

### ASIC: synthesizable

961 e.g. / 20.0 ns  
2 sub-s, 2 comp-s

FPGA: non-synthesizable  
wait statements in loops :(  
explicit FSM needed :( :(

### Possible trade-offs

- functional unit sharing
- universal functional units
- out-of-order execution



## GCD example – behavioral FSM

```

process begin -- gcd-bfsm.vhdl
  wait on clk until clk='1';
  case state is
  -- Wait for the new input data
  when S_wait =>
    if rst='0' then
      x<=xi; y<=yi; rdy<='0'; state<=S_start;
    end if;
  -- Calculate
  when S_start =>
    if x /= y then
      if x < y then y <= y - x;
      else x <= x - y; end if;
      state<=S_start;
    else
      xo<=x; rdy<='1'; state<=S_ready;
    end if;
  -- Ready
  when S_ready => state<=S_wait;
  end case;
end process;

```

ASIC: synthesizable  
 911 e.g. / 19.4 ns  
 2 sub-s, 2 comp-s

FPGA: synthesizable  
 108 SLC / 9.9 ns  
 2 sub-s, 2 comp-s

Can it be made better?

Again the possible trade-offs

- functional unit sharing  
     one operation per clock step
- universal functional units  
      $A < B \implies A - B < 0$  /  $A = B \implies A - B = 0$
- out-of-order execution  
     subtracting first then deciding



## GCD example – universal functional units?

- $A < B \implies A - B < 0$  /  $A = B \implies A - B = 0$

```

-- Three operations:
-- subtraction, and
-- comparisons not-equal &
-- less-than
xo <= xi - yi;

ne <= '1' when xi /= yi else '0';

lt <= '1' when xi < yi else '0';

```

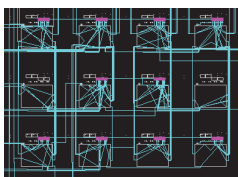
```

-- ALU - subtracting and then comparing
x_out <= xi - yi;    xo <= x_out;
process (x_out)
  variable or_tmp: unsigned(15 downto 0);
begin
  or_tmp(15) := x_out(15);
  for i in 14 downto 0 loop
    or_tmp(i) := or_tmp(i+1) or x_out(i);
  end loop;
  ne <= to_bit(or_tmp(0));
end process;
lt <= to_bit(x_out(15));

```

ASIC: 209 e.g. / 19.9 ns  
 FPGA: 20 SLC / 12.1 ns  
three adder chains!

ASIC: 148 e.g. / 21.8 ns / FPGA: 12 SLC / 14.6 ns





## GCD example – design space exploration

- Different solutions – <http://mini.pld.ttu.ee/~lrv/gcd/gcd.html>
- gcd-bhv.vhdl – pure behavioral description, non-synthesizable
- gcd-bhvc.vhdl – fully clocked behavioral style, some synthesis tools can handle
- gcd-bfsm.vhdl – so called behavioral FSM (explicit FSM & behavioral data-path), synthesizable (but how efficient it is?)
- gcd-rtl1.vhdl – single ALU, 3 clock cycles per iteration – 1) “not equal?”, 2) “less than?”, 3) subtract
- gcd-rtl2.vhdl – single ALU, 2 clock cycles per iteration – 1) “not equal?” and “less than?”, 2) subtract
- gcd-rtl3.vhdl – comparator (less than) controls subtraction, 1 clock cycle per iteration – small but slow (sequential) data-path
- gcd-rtl4.vhdl – out-of-order execution – both subtractions are calculated first then the decision is made (one subtracter compares for “less than”, another for “not equal”)
- gcd-rtl5.vhdl – out-of-order execution – both subtractions are calculated first then the decision is made (one subtracter compares for “less than” but separate “not equal”)



## GCD example – single ALU, 2 clock cycles per iteration

### *gcd-rtl2.vhdl*

```

-- Next state function of the FSM
process (state, rst, alu_ne, alu_lt) begin
  ena_x <= '0';   ena_y <= '0';   ena_r <= '0';
  set_rdy <= '0'; xi_yi_sel <= '0'; sub_y_x <= '0';
  next_state <= state;
  case state is
    when S_wait => -- Wait for the new input data
      if rst='0' then
        xi_yi_sel <= '1'; ena_x <= '1'; ena_y <= '1';
        next_state <= S_start;
      end if;
    when S_start => -- Loop: ready?
      if alu_ne='1' then
        if alu_lt='1' then next_state <= S_sub_y_x;
        else next_state <= S_sub_x_y; end if;
      else next_state <= S_ready;
      end if;
    when S_sub_y_x => -- Loop: y-x
      ena_y <= '1'; sub_y_x <= '1';
      next_state <= S_start;
    when S_sub_x_y => -- Loop: x-y
      ena_x <= '1'; sub_y_x <= '0';
      next_state <= S_start;
    when S_ready => -- Ready
      ena_r <= '1'; set_rdy <= '1';
      next_state <= S_wait;
  end case;
end process;

-- ALU: subtract / less-than / not-equal
alu_o <= alu_1 - alu_2;
alu_lt <= to_bit(alu_o(15));
process (alu_o)
  variable or_tmp: unsigned(15 downto 0);
begin
  or_tmp(15) := alu_o(15);
  for i in 14 downto 0 loop
    or_tmp(i) := or_tmp(i+1) or alu_o(i);
  end loop;
  alu_ne <= to_bit(or_tmp(0));
end process;

-- Multiplexers
x_i <= xi when xi_yi_sel='1' else alu_o;
y_i <= yi when xi_yi_sel='1' else alu_o;
alu_1 <= y when sub_y_x='1' else x;
alu_2 <= x when sub_y_x='1' else y;

-- Registers
process begin
  wait on clk until clk='1';
  state <= next_state;
  if ena_x='1' then x <= x_i; end if;
  if ena_y='1' then y <= y_i; end if;
  if ena_r='1' then xo <= x; end if;
  rdy <= set_rdy;
end process;

```



## GCD example – out-of-order execution (2 sub-s)

*gcd-rtl5.vhdl*

```

-- Next state function of the FSM
process (state, rst, alu_ne) begin
ena_xy <= '0';    ena_r <= '0';
set_rdy <= '0';  xi_yi_sel <= '0';
next_state <= state;
case state is
-- Wait for the new input data
when S_wait =>
if rst='0' then
xi_yi_sel <= '1';  ena_xy <= '1';
next_state <= S_start;
end if;
-- Calculate
when S_start =>
if alu_ne='1' then
ena_xy <= '1';
next_state <= S_start;
else
ena_r <= '1';    set_rdy <= '1';
next_state <= S_ready;
end if;
-- Ready
when S_ready =>
ena_r <= '1';    set_rdy <= '1';
next_state <= S_wait;
end case;
end process;

-- Subtractor (x-y) / comparator (x<y)
alu_o1 <= x - y;
sub_y_x <= '1' when alu_o1(alu_o1'high)='1' else '0';

-- Subtractor (y-x)
alu_o2 <= y - x;

-- Comparator (y/=x)
alu_ne <= '1' when x /= y else '0';

-- Multiplexers
x_i <= xi when xi_yi_sel='1' else alu_o1;
y_i <= yi when xi_yi_sel='1' else alu_o2;
ena_x <= '1' when (sub_y_x='0' and ena_xy='1') or
xi_yi_sel='1' else '0';
ena_y <= '1' when (sub_y_x='1' and ena_xy='1') or
xi_yi_sel='1' else '0';

-- Registers
process begin
wait on clk until clk='1';
state <= next_state;
if ena_x='1' then x <= x_i;  end if;
if ena_y='1' then y <= y_i;  end if;
if ena_r='1' then xo <= x;  end if;
rdy <= set_rdy;
end process;

```



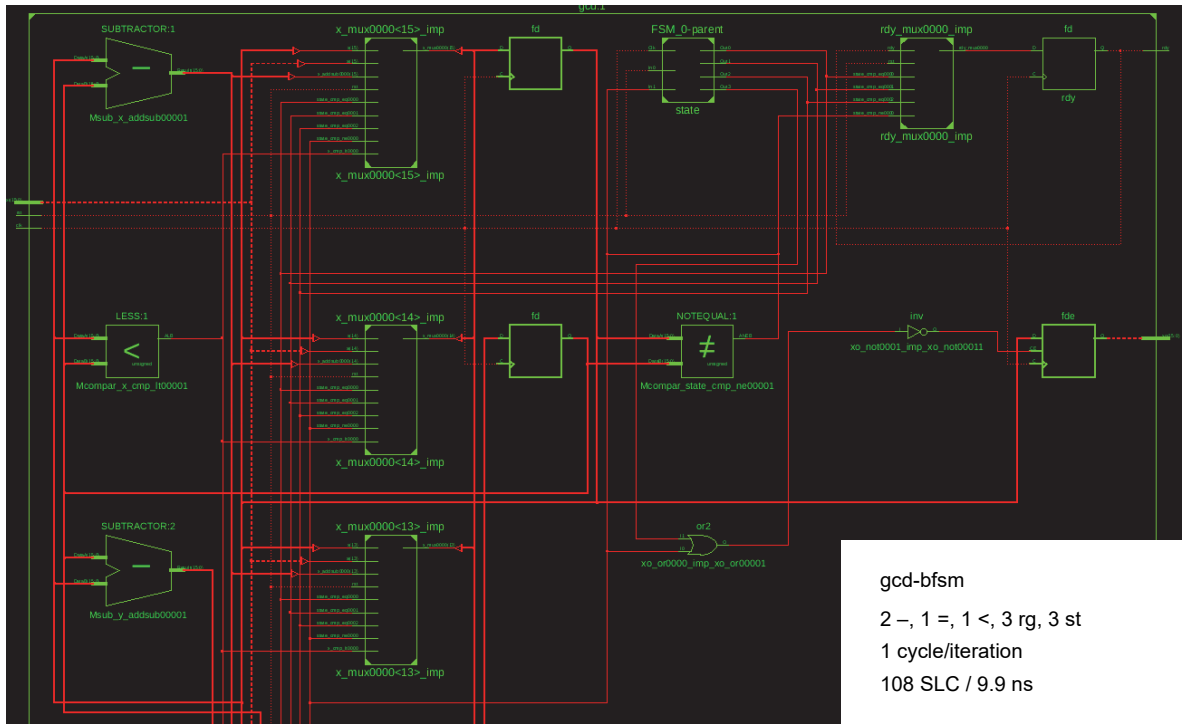
## GCD example – synthesis results

Technology	FPGA				ASIC			
	50 MHz		100 MHz		50 MHz		25 MHz	
Constraint <sup>1)</sup>	[SLC]	[ns]	[SLC]	[ns]	[e.g.]	[ns]	[e.g.]	[ns]
<i>gcd-bhv</i> <sup>2)</sup>	93	17.3	-	-	1141	20.0	-	-
gcd-bhvc	-	-	-	-	961	20.0	977	31.1
gcd-bfsm	108	9.9	108	9.4	911	19.4	984	30.8
gcd-rtl1	50	10.8	50	9.7	986	19.8	883	32.4
gcd-rtl2	48	10.8	48	10.0	931	19.9	882	32.3
gcd-rtl3	58	17.0	58	14.6	1134	20.0	928	40.0
gcd-rtl4	78	12.6	78	9.0	976	19.9	928	29.0
gcd-rtl5	58	8.0	58	7.6	915	20.0	932	26.9

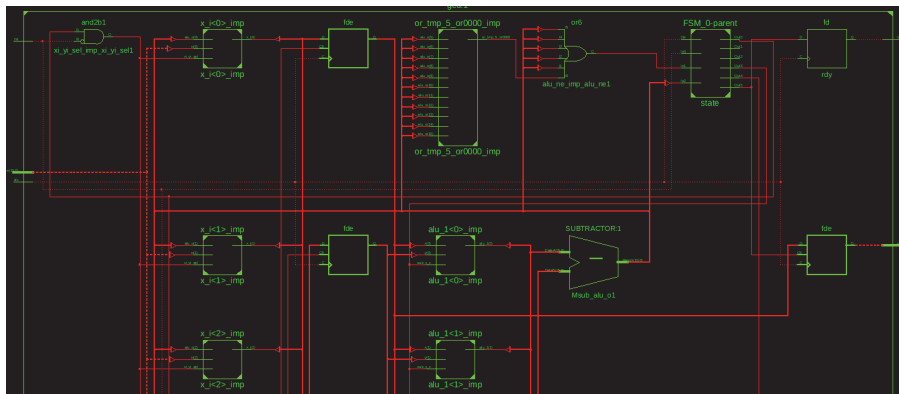
<sup>1)</sup> Clock period was the only constraint

<sup>2)</sup> gcd-bhv was synthesized using the help of prototype HLS tool xTractor

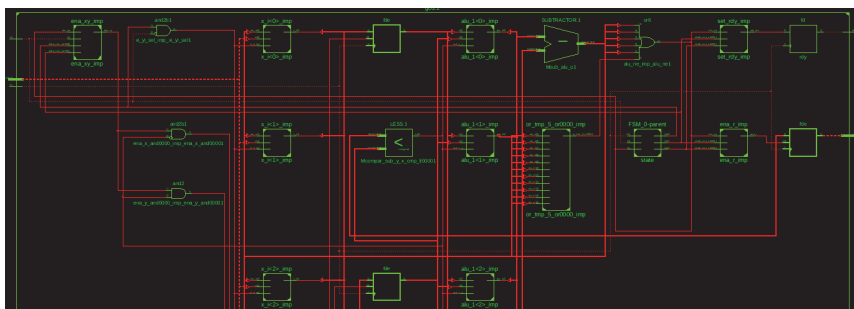
## GCD example – why such differences?



## GCD example – why such differences?



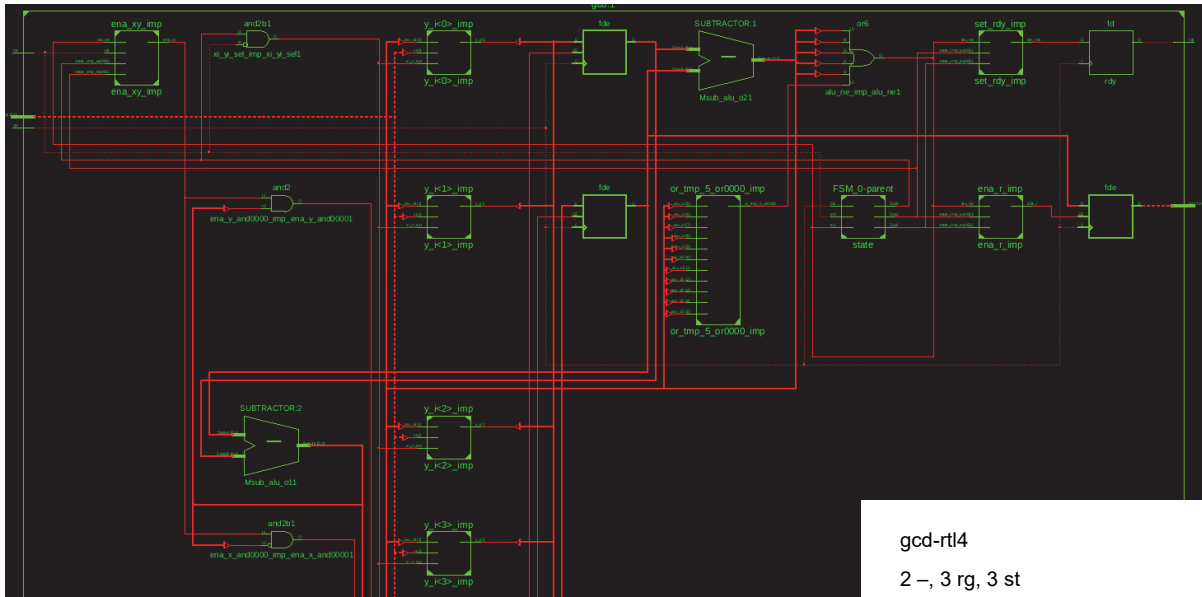
gcd-rtl1  
 1 ALU, 3 rg, 6 st  
 3 cycles/iteration  
 50 SLC / 10.8 ns



gcd-rtl2  
 1 ALU, 3 rg, 5 st  
 2 cycles/iteration  
 48 SLC / 10.8 ns

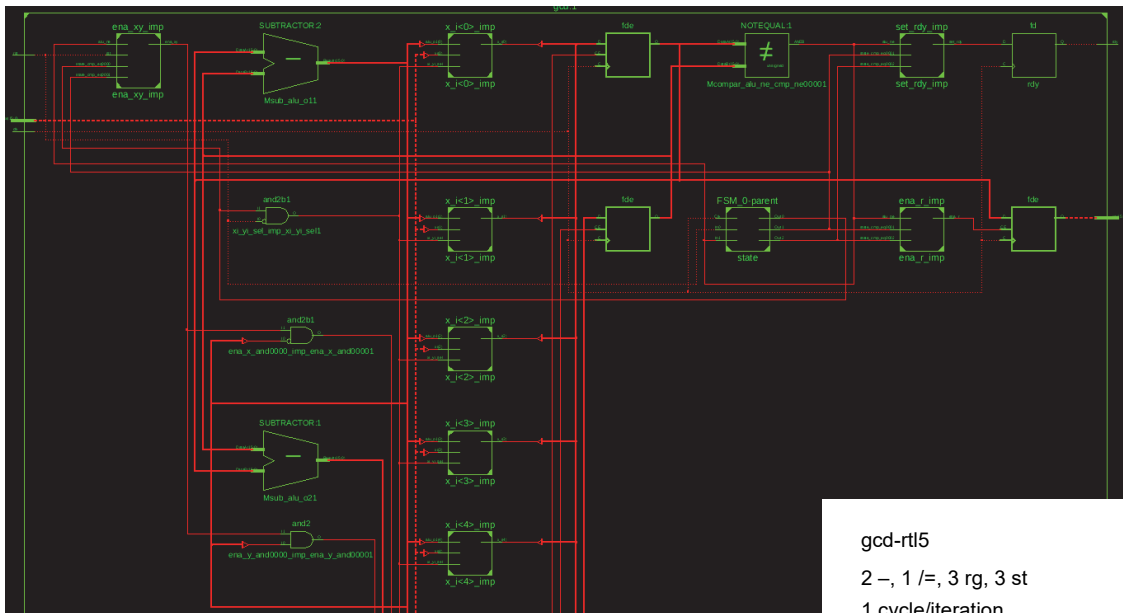
gcd-rtl3  
 1 -, 1 <, 3rg, 3st  
 1 cycle/iteration  
 58 SLC / 17.0 ns

## GCD example – why such differences?



gcd-rtl4  
 2 -, 3 rg, 3 st  
 1 cycle/iteration  
 78 SLC / 12.6 ns

## GCD example – why such differences?



gcd-rtl5  
 2 -, 1 /=, 3 rg, 3 st  
 1 cycle/iteration  
 58 SLC / 8.0 ns